

# Linear Convergence of Evolution Strategies with Derandomized Sampling

## Beyond Quasi-Convex Functions

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# Introduction

The aim of this presentation

Study convergence rate of a simple pattern search method

# Introduction

## Convergence of evolutionary algorithms

Proofs are almost always for (quasi-)convex objective functions

## Non quasi-convex objective functions

- ▶ some proofs for convergence (asymptotically the optimum is found)
- ▶ few for linear convergence (precision  $O(e^{-\Omega(n)})$  after  $n$  iterations)
  - ▶ in discrete space
  - ▶ not in continuous space

# Introduction

## Our contribution

- ▶ Prove the linear convergence of an algorithm
  - ▶ on non quasi-convex functions
  - ▶ on continuous domains
- ▶ Under some assumptions about
  - ▶ the sampling performed by the algorithm
  - ▶ the “conditioning” of the objective function
  - ▶ the unicity of the optimum

# Overview

Framework

Theorem and proof

Application to quadratic functions

Conclusion

# Framework

# Algorithm

Initialize  $\mathbf{x} \in \mathbb{R}^d$

Parameters  $k \in \mathbb{N}^*$ ,  $\delta_1, \dots, \delta_k \in \mathbb{R}^d$ ,  $\sigma \in \mathbb{R}_+^*$ ,  $k_1 \in \mathbb{N}^*$ ,  $k_2 \in \mathbb{N}^*$

**for**  $t = 1, 2, 3, \dots$  **do**

// mutations

For  $i \in [[1, k]]$ ,  $\mathbf{x}_i \leftarrow \mathbf{x} + \sigma \delta_i$

// useful auxiliary variables

$n \leftarrow$  number of  $\mathbf{x}_i$  such that  $f(\mathbf{x}_i) < f(\mathbf{x})$

$\mathbf{x}' \leftarrow \mathbf{x}_i$  with  $i \in [[1, k]]$  such that  $f(\mathbf{x}_i)$  is minimum

...

**end for**

# Algorithm

Initialize  $\mathbf{x} \in \mathbb{R}^d$

Parameters  $k \in \mathbb{N}^*$ ,  $\delta_1, \dots, \delta_k \in \mathbb{R}^d$ ,  $\sigma \in \mathbb{R}_+^*$ ,  $k_1 \in \mathbb{N}^*$ ,  $k_2 \in \mathbb{N}^*$

**for**  $t = 1, 2, 3, \dots$  **do**

...

// step-size adaptation

**if**  $n \leq k_1$  **then**

$\sigma \leftarrow \sigma/2$

**end if**

**if**  $n \geq k_2$  **then**

$\sigma \leftarrow 2\sigma$

**end if**

...

**end for**



# Algorithm

```

Initialize  $\mathbf{x} \in \mathbb{R}^d$ 
Parameters  $k \in \mathbb{N}^*$ ,  $\delta_1, \dots, \delta_k \in \mathbb{R}^d$ ,  $\sigma \in \mathbb{R}_+^*$ ,  $k_1 \in \mathbb{N}^*$ ,  $k_2 \in \mathbb{N}^*$ 
for  $t = 1, 2, 3, \dots$  do

    ...

    // win: accepted mutation
    if  $k_1 < n < k_2$  then
         $\mathbf{x} \leftarrow \mathbf{x}'$ 
    end if

end for

```

# Assumptions

## Objective function

The objective function  $f$  is unimodal

The considered algorithms are invariant by translation or composition with increasing functions, therefore we can state that

- ▶  $\mathbf{x}^* = \mathbf{0}$  is the optimum
- ▶  $f(\mathbf{x}^*) = 0$

# Assumptions

## Conditioning of $f$

Conditioning of  $f$ :  $\exists K' > 0, \exists K'' > 0$  s.t.  $\forall \mathbf{x} \in \mathbb{R}^d$

$$K' \|\mathbf{x}\| \leq f(\mathbf{x}) \leq K'' \|\mathbf{x}\|$$

This assumption is not so strong as a constraint and in fact, quadratic positive definite forms with bounded condition number are covered

# Assumptions

## Deterministic sampling of the algorithm

The sampling of the algorithm is deterministic (like in pattern search methods)

- ▶ mutation vectors  $\delta_i$  are constant
- ▶ the evolution of the step size parameter  $\sigma$  is deterministic

# Assumptions

## Regular sampling of the algorithm

We assume  $\exists b, b', c', c, \eta$  s.t.

$$0 < b < b' \leq 2b' \leq c' \leq c, \quad 0 < \eta < 1, \quad \forall \mathbf{x} \in \mathbb{R}^d$$

$$\sigma \geq b^{-1} \|\mathbf{x}\| \Rightarrow n \leq k_1 \quad (\sigma \text{ too large}) \quad (1)$$

$$\sigma \leq b'^{-1} \|\mathbf{x}\| \Rightarrow n > k_1 \quad (\sigma \text{ small enough}) \quad (2)$$

$$\sigma \geq c'^{-1} \|\mathbf{x}\| \Rightarrow n < k_2 \quad (\sigma \text{ large enough}) \quad (3)$$

$$\sigma \leq c^{-1} \|\mathbf{x}\| \Rightarrow n \geq k_2 \quad (\sigma \text{ too small}) \quad (4)$$

$$b'^{-1} \|\mathbf{x}\| \leq \sigma \leq c'^{-1} \|\mathbf{x}\| \Rightarrow \exists i \in [[1, k]]; f(\mathbf{x}_i) \leq \eta f(\mathbf{x}) \quad (5)$$

with  $n := \#\{i \in [[1, k]]; f(\mathbf{x} + \sigma \delta_i) < f(\mathbf{x})\}$

# Theorem

## Preliminary work

Define  $l = \ln\left(\frac{\|\mathbf{x}\|}{\sigma}\right)$ . Eqs. 1-5 can be rephrased as follows:

$$l \leq \ln(b) \Rightarrow n \leq k_1 \quad (\sigma \text{ too large}) \quad (6)$$

$$l \geq \ln(b') \Rightarrow n > k_1 \quad (\sigma \text{ small enough}) \quad (7)$$

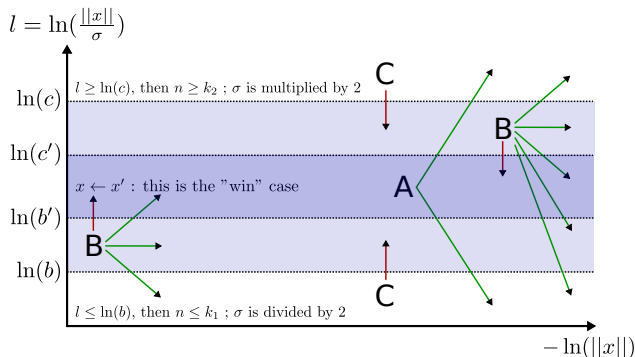
$$l \leq \ln(c') \Rightarrow n < k_2 \quad (\sigma \text{ large enough}) \quad (8)$$

$$l \geq \ln(c) \Rightarrow n \geq k_2 \quad (\sigma \text{ too small}) \quad (9)$$

$$\ln(b') \leq l \leq \ln(c') \Rightarrow \exists i \in [[1, k]]; f(\mathbf{x} + \sigma\delta_i) \leq \eta f(\mathbf{x}) \quad (10)$$

with  $n := \#\{i \in [[1, k]]; f(\mathbf{x} + \sigma\delta_i) < f(\mathbf{x})\}$

# Preliminary work





# Preliminary work

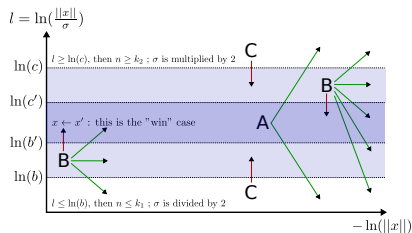
## Forced increase

### Forced increase

if  $l \leq \ln(b)$ , then

- ▶  $n \leq k_1$
- ▶  $\sigma$  is divided by 2
- ▶  $l$  is increased by  $\ln(2)$  (Eq. 6)

This is a case C at the bottom in the figure



# Preliminary work

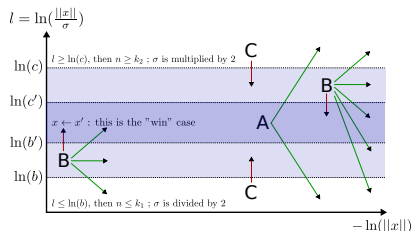
## Forced decrease

### Forced decrease

if  $l \geq \ln(c)$ , then

- ▶  $n \geq k_2$
- ▶  $\sigma$  is multiplied by 2
- ▶  $l$  is decreased by  $\ln(2)$  (Eq. 9)

This is a case C at the top in the figure



# Preliminary work

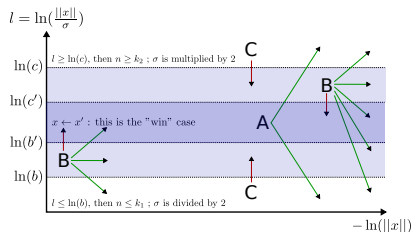
## Forced win

### Forced win

if  $\ln(b') \leq l \leq \ln(c')$ , then

- ▶ this is the “sure win” case (Eq. 10)
- ▶  $\mathbf{x} \leftarrow \mathbf{x}'$  ( $\mathbf{x}'$  is the best  $\mathbf{x}_j$ )
- ▶  $l$  can be
  - ▶ increased (at most by  $\max_i \|\delta_i\|$ ) or
  - ▶ decreased (by  $\Delta = \ln\left(\frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|}\right)$ )

This is a case A in the figure



# Preliminary work

## Uncertain outcome

### Uncertain outcome

if  $\ln(b) \leq l \leq \ln(b')$  or  $\ln(c) \leq l \leq \ln(c')$ , then

- ▶ the iteration can be an improvement or not
- ▶ if not, the point is moved towards case A with steps of  $\ln(2)$
- ▶ if there's no “win” case, then in the mean time  $l$  will arrive between  $\ln(b')$  and  $\ln(c')$ , where a win is ensured

This is a case B in the previous figure

# Theorem

There exists a constant  $K$ , depending on  $\eta, K', K'', \max_i \|\delta_i\|$  only such that for index  $t$  large enough

$$\frac{\ln(\|\mathbf{X}_t\|)}{t} \leq K < 0$$

where  $\mathbf{X}_t$  is the tested solution  $\mathbf{x}$  at iteration  $t$

# Proof

## Step 1

### Showing that there are infinitely many wins

1.  $l$  is increased or decreased when it is too low or too high
  - ▶ the algorithm eventually brings  $l$  to the “win” range
2.  $l$  can be increased or decreased at most by  $\ln(2)$  and  $b' \leq 2b' \leq c'$ 
  - ▶ the algorithm can not jump over the “win” range

This ensures that infinitely often we have a “win” :  $\mathbf{x} \leftarrow \mathbf{x}'$

# Proof

## Step 2

Showing that “wins” are big enough

“Win” case implies

- ▶  $f(\mathbf{x}') \leq \eta f(\mathbf{x})$
- ▶  $f(\mathbf{x}') \leq K'' \|\mathbf{x}'\| \leq \frac{K''}{K'} \frac{\|\mathbf{x}'\|}{\|\mathbf{x}\|} f(\mathbf{x})$

so that  $\ln(f(\mathbf{x}))$  is decreased by at least

$$\max \left( \ln \left( \frac{1}{\eta} \right), \ln \left( \frac{K'}{K''} \right) + \ln \left( \frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|} \right) \right) \quad (11)$$

# Proof

## Step 2

Showing that the number of steps between two “wins” is low enough

After a “win”, the number of iterations to the next “win” is

- ▶ at most  $z = 1 + \ln\left(\frac{c}{b}\right) \frac{\Delta}{\ln(2)}$  if  $l' \leq \ln(b')$
- ▶ at most  $z = 1 + \frac{\max_i \|\delta_i\|}{\ln(2)}$  if  $l' \geq \ln(c')$
- ▶ less than both cases above otherwise

with  $l' = \ln\left(\frac{\|\mathbf{x}'\|}{\sigma}\right)$  and  $\Delta = \ln\left(\frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|}\right)$



# Proof

## Step 2

### Progress rate

Eq. 11 divided by  $z$  is lower bounded by some positive constant

$$\begin{aligned} \text{ProgressRate} &= \text{Eq. 11 divided by } z \\ &= \frac{\max\left(\ln\left(\frac{1}{\eta}\right), \ln\left(\frac{K'}{K''}\right) + \Delta\right)}{\min\left(1 + \ln\left(\frac{c}{b}\right) \frac{\Delta}{\ln(2)}, 1 + \frac{\max_j \|\delta_j\|}{\ln(2)}\right)} \end{aligned}$$

# Proof

## Step 3

### Summing iterations

Hence if  $t > n_0$ ,

$$\ln(f(\mathbf{X}_t)) \leq \ln(f(\mathbf{X}_1)) - (t - n_0) \times \sum_i \frac{\max\left(\ln\left(\frac{1}{\eta}\right), \ln\left(\frac{K'}{K''}\right) + \Delta\right)}{\min\left(1 + \ln\left(\frac{c}{b}\right) \frac{\Delta}{\ln(2)}, 1 + \frac{\max_j \|\delta_j\|}{\ln(2)}\right)} \quad \square$$

where :

- ▶ summation is for  $i$  index of an iteration  $t$  with a “win”
- ▶  $n_0$  is the number of initial iterations before a “win”

# Application to quadratic functions

## Application to quadratic functions

Considered family of objective functions

$f$  is quadratic positive definite objective functions such that

$$\frac{\max \text{EigenValue}(\text{Hessian}(f))}{\min \text{EigenValue}(\text{Hessian}(f))} < C_{\max} < \infty$$

Consider  $Q$  a positive definite quadratic form with optimum in 0

We work on  $\mathbf{x} \mapsto \sqrt{Q(\mathbf{x} - \mathbf{x}^*)}$  instead of  $\mathbf{x} \mapsto Q(\mathbf{x} - \mathbf{x}^*)$  so that the first assumption is verified:

$$K' \|\mathbf{x}\| \leq f(\mathbf{x}) \leq K'' \|\mathbf{x}\|$$

$$\forall \mathbf{x} \in \mathbb{R}^d, \exists K' > 0, \exists K'' > 0$$

# Application to quadratic functions

## Assumptions to verify (reminder)

We try to prove that  $f$  respects the following assumptions

$\exists b, b', c', c, \eta$  s.t.  $0 < b < b' \leq 2b' \leq c' \leq c$ ,  $0 < \eta < 1$ ,  $\forall \mathbf{x} \in \mathbb{R}^d$

$$\sigma \geq b^{-1} \|\mathbf{x}\| \Rightarrow n \leq k_1 \quad (\sigma \text{ too large})$$

$$\sigma \leq b'^{-1} \|\mathbf{x}\| \Rightarrow n > k_1 \quad (\sigma \text{ small enough})$$

$$\sigma \geq c'^{-1} \|\mathbf{x}\| \Rightarrow n < k_2 \quad (\sigma \text{ large enough})$$

$$\sigma \leq c^{-1} \|\mathbf{x}\| \Rightarrow n \geq k_2 \quad (\sigma \text{ too small})$$

$$b'^{-1} \|\mathbf{x}\| \leq \sigma \leq c'^{-1} \|\mathbf{x}\| \Rightarrow \exists i \in [[1, k]]; f(\mathbf{x}_i) \leq \eta f(\mathbf{x})$$

with  $n := \#\{i \in [[1, k]]; f(\mathbf{x} + \sigma \delta_i) < f(\mathbf{x})\}$

# Application to quadratic functions

We note:

- ▶  $p = p_{\mathbf{x},\sigma,f}$  the probability that  $\mathbf{x} + \sigma\delta_i$  is in  $E = f^{-1}([0, f(\mathbf{x})])$
- ▶  $\hat{p} = \hat{p}_{\mathbf{x},\sigma,f}$  the frequency  $\frac{1}{k} \sum_{i=1}^k \mathbf{1}_{\mathbf{x}+\sigma\delta_i \in E}$

The previous assumptions essentially mean that frequencies are close to expectations for

- ▶  $\mathbf{x} + \sigma\delta_i \in f^{-1}([0, f(\mathbf{x})])$
- ▶  $\mathbf{x} + \sigma\delta_i \in f^{-1}([0, \eta f(\mathbf{x})])$

uniformly in  $\mathbf{x}, \sigma, f$ .

# Corollary

## Application of the main theorem to quadratic forms

Assume that the  $\delta_i$  are uniformly randomly drawn in the unit ball  $B(0, 1)$ .

Assume that  $F$  is the set of quadratic functions with minimum in 0 ( $f(0) = 0$ ) as defined before.

Then, almost surely on the sequence  $\delta_1, \delta_2, \dots, \delta_k$ , for  $k$  large enough and some parameters  $k_1$  and  $k_2$  of our evolutionary algorithm, then assumptions hold, and therefore for some  $K < 0$ , for all  $t > 0$ ,

$$\frac{\ln(\|\mathbf{X}_t\|)}{t} \leq K$$

where  $\mathbf{X}_t$  is the tested solution  $\mathbf{x}$  at iteration  $t$

# Proof

## Step 1

### Using VC-dimension for approximating expectations by frequencies

The finiteness of the VC-dimension of quadratic forms state that for all  $\epsilon > 0$ , almost surely in  $\delta_1, \delta_2, \dots, \delta_k$ , for all  $\delta > 0$  and  $k$  sufficiently large, with probability at least  $1 - \delta$ ,

$$\sup_{\mathbf{x}, f, \sigma > 0} |\hat{p}_{\mathbf{x}, \sigma, f} - p_{\mathbf{x}, \sigma, f}| \leq \epsilon/2$$

where  $\mathbf{x}$  ranges over the domain,  $f$  ranges over  $F$



# Proof

## Step 2

Showing that small  $\sigma$  leads to high acceptance rate and high  $\sigma$  leads to small acceptance rate

Thanks to the bounded conditioning, there exists  $\epsilon > 0$  s.t.

$$s' < \frac{1}{2}s$$

$$\text{with } s = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } p \geq \frac{\epsilon}{2} \right\}$$

$$\text{and } s' = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } p < \frac{1}{2} - \frac{\epsilon}{2} \right\}$$

Indeed  $s' \rightarrow 0$  and  $s \rightarrow \infty$  as  $\epsilon \rightarrow 0$ .

# Proof

## Step 2

Showing that small  $\sigma$  leads to high acceptance rate and high  $\sigma$  leads to small acceptance rate

The previous equations provide  $k_1$ ,  $k_2$ ,  $c'$  and  $b'$

$$\frac{1}{b'} = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} \geq \epsilon \right\}$$

$$\frac{1}{c'} = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} < \frac{1}{2} - \epsilon \right\}$$

$$k_1 = \lfloor \epsilon k \rfloor$$

$$k_2 = \left\lceil \left( \frac{1}{2} - \epsilon \right) k \right\rceil$$

Equations above imply  $c' \geq 2b'$

# Proof

## Step 3

Showing that  $k$  large enough and  $\sigma$  well chosen leads to at least one mutation with significant improvement

Similarly,  $k$  large enough yield

$$b^{-1} = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} > k_1/k \right\}$$

$$c^{-1} = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} < k_2/k \right\}$$

which provide Eqs. 4 and 1 with  $b < c$ .

Eqs. 1-4 then imply  $b < b'$  and  $c' < c$ .

# Proof

## Step 3

Showing that  $k$  large enough and  $\sigma$  well chosen leads to at least one mutation with significant improvement

We now have to ensure the last assumption:

$$b'^{-1} \|\mathbf{x}\| \leq \sigma \leq c'^{-1} \|\mathbf{x}\| \Rightarrow \exists i \in [[1, k]]; f(\mathbf{x} + \sigma \delta_i) \leq \eta f(\mathbf{x})$$

For now on, we note:

- ▶  $q = q_{\mathbf{x}, \sigma, f}$  the probability that  $\mathbf{x} + \sigma \delta_i$  is in  $E' = f^{-1}([0, \eta f(\mathbf{x})])$
- ▶  $\hat{q} = \hat{q}_{\mathbf{x}, \sigma, f}$  the frequency  $\frac{1}{k} \sum_{i=1}^k \mathbf{1}_{\mathbf{x} + \sigma \delta_i \in E'}$

# Proof

## Step 3

Showing that  $k$  large enough and  $\sigma$  well chosen leads to at least one mutation with significant improvement

Lets us assume

$$b^{-1} \leq \frac{\sigma}{\|\mathbf{x}\|} \leq c^{-1}$$

this implies  $q > \epsilon_0$  for some  $\epsilon_0 > 0$

For  $k$  sufficiently large for ensuring  $\sup_{\sigma, \mathbf{x}, f} |q_{\mathbf{x}, \sigma, f} - \hat{q}_{\mathbf{x}, \sigma, f}| \leq \epsilon_0/2$ , by VC-dimension, we get  $q' \geq \epsilon_0/2 > 0$

This implies that at least one  $\delta_i$  verifies  $\mathbf{x} + \delta_i \in E'$ .

This is the last assumption.

# Proof

## Step 4

### Concluding

We have shown our assumptions for square roots of positive definite quadratic normal forms with bounded conditioning. Therefore, the main theorem can be applied and leads to

$$\frac{\ln(\|\mathbf{X}_t\|)}{t} \leq K < 0$$

# Conclusion

# Conclusion

This is the first proof of linear convergence of an evolutionary algorithm in continuous domains on non quasi-convex functions.

Even the application to quadratic positive definite forms is new.



## Future work

- ▶ Evaluate the convergence rate as a function of condition numbers
- ▶ Extend results to randomized algorithms

# Thank you for your attention

Questions ?

# Theorem's proof

## Step 2

$$(11) \iff \ln \left( \frac{f(\mathbf{x})}{f(\mathbf{x}')} \right) \geq \max \left( \ln \left( \frac{1}{\eta} \right), \ln \left( \frac{K'}{K''} \right) + \ln \left( \frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|} \right) \right)$$

$$f(\mathbf{x}') \leq \eta f(\mathbf{x}) \iff \frac{1}{\eta} f(\mathbf{x}') \leq f(\mathbf{x})$$

$$\iff \frac{1}{\eta} \leq \frac{f(\mathbf{x})}{f(\mathbf{x}')}$$

$$\iff \ln \left( \frac{f(\mathbf{x})}{f(\mathbf{x}')} \right) \geq \ln \left( \frac{1}{\eta} \right)$$

$$f(\mathbf{x}') \leq \frac{K''}{K'} \frac{\|\mathbf{x}'\|}{\|\mathbf{x}\|} f(\mathbf{x}) \iff f(\mathbf{x}') \frac{K'}{K''} \frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|} \leq f(\mathbf{x})$$

$$\iff \frac{K'}{K''} \frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|} \leq \frac{f(\mathbf{x})}{f(\mathbf{x}')}$$

$$\iff \ln \left( \frac{K'}{K''} \frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|} \right) \leq \ln \left( \frac{f(\mathbf{x})}{f(\mathbf{x}')} \right)$$

$$\iff \ln \left( \frac{f(\mathbf{x})}{f(\mathbf{x}')} \right) \geq \ln \left( \frac{K'}{K''} \right) + \ln \left( \frac{\|\mathbf{x}\|}{\|\mathbf{x}'\|} \right)$$

# Theorem's proof

## Step 2

if  $l' \geq \ln(c')$  then

- ▶  $l < \ln(c)$  (otherwise it couldn't be a "win")

$$\begin{aligned}
 l' - \ln(c') &\leq \ln(c + \max_i \|\delta_i\|) - \ln(c') \\
 &\leq \ln(c) + \ln(1 + \max_i \|\delta_i\|/c) - \ln(c') \\
 &\leq \ln(c/c') + \max_i \|\delta_i\|/c \\
 &\leq \max_i \|\delta_i\|/c
 \end{aligned}$$

$$\begin{aligned}
 l' - \ln(c') = \ln\left(\frac{\|x + \sigma\delta_i\|}{\sigma}\right) - \ln(c') &\leq \ln\left(\frac{\|x\|}{\sigma} + \delta_i\right) - \ln(c') \\
 &\leq \ln\left(c * \left(1 + \frac{\delta_i}{c}\right)\right) - \ln(c') \\
 &\leq \ln(c) + \ln\left(1 + \max_i \|\delta_i\|/c\right) - \ln(c') \\
 &\leq \ln(c/c') + \max_i \|\delta_i\|/c \\
 &\leq \max_i \|\delta_i\|/c
 \end{aligned}$$

# Theorem's Proof

## Step 3

### Summing iterations

Hence if  $t > n_0$ ,

$$\ln(f(\mathbf{X}_t)) \leq \ln(f(\mathbf{X}_1)) - (t - n_0) \times \sum_i \frac{\max\left(\ln\left(\frac{1}{\eta}\right), \ln\left(\frac{K'}{K''}\right) + \Delta\right)}{\min\left(1 + \ln\left(\frac{\epsilon}{b}\right) \frac{\Delta}{\ln(2)}, 1 + \ln\left(\frac{\epsilon}{b}\right) \frac{\max_j \ln(\|\delta_j\|)}{\ln(2)}\right)}$$

$$\Leftrightarrow \ln(f(\mathbf{X}_t)) - \ln(f(\mathbf{X}_1)) \leq -(t - n_0) \times C$$

$$\Leftrightarrow \frac{\ln\left(\frac{f(\mathbf{X}_t)}{f(\mathbf{X}_1)}\right)}{t - n_0} \leq -C$$

$$\Rightarrow \frac{\ln(\|\mathbf{X}_t\|)}{t} \leq K < 0 \quad (\text{theorem})$$

with  $C$  a positive constant.

## Corollary's proof

### Step 2 (1/4)

Step 2: showing that  $\sigma$  small leads to high acceptance rate and  $\sigma$  high leads to small acceptance rate.

Thanks to the bounded conditioning, there exists  $\epsilon > 0$  s.t.

$$s' < \frac{1}{2}s \quad (12)$$

$$\text{with } s = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } p \geq \frac{\epsilon}{2} \right\}$$

$$\text{and } s' = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } p < \frac{1}{2} - \frac{\epsilon}{2} \right\}$$

because  $s' \rightarrow 0$  and  $s \rightarrow \infty$  as  $\epsilon \rightarrow 0$ .

# Corollary's proof

## Step 2 (2/4)

Notes

$$\hat{s} = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} \geq \epsilon \right\}$$

$$\hat{s}' = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} < \frac{1}{2} - \epsilon \right\}$$

Then

$$\sup_{\mathbf{x}, f, \sigma > 0} |\hat{p} - p| \leq \epsilon/2$$

implies

$$\frac{1}{2}\hat{s} \geq \frac{1}{2}s \quad \text{and} \quad s' \geq \hat{s}'$$

# Corollary's proof

## Step 2 (3/4)

So

$$\hat{s}' \leq s' < \frac{1}{2}s \leq \frac{1}{2}\hat{s}$$

and

$$\hat{s}' \leq \frac{1}{2}\hat{s}$$



# Corollary's proof

## Step 2 (4/4)

This provides  $k_1$ ,  $k_2$ ,  $c'$  and  $b'$  as follows for Eqs. 3 and 2:

$$\begin{aligned} \frac{1}{b'} &= \hat{s} = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} \geq \epsilon \right\} \\ \frac{1}{c'} &= \hat{s}' = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} < \frac{1}{2} - \epsilon \right\} \\ k_1 &= \lfloor \epsilon k \rfloor \\ k_2 &= \left\lceil \left( \frac{1}{2} - \epsilon \right) k \right\rceil \end{aligned}$$

Eqs. above imply  $c' \geq 2b'$

# Corollary's proof

## Step 3

$k$  large enough yield

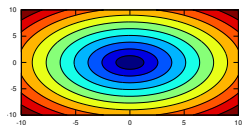
$$b^{-1} = \sup \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} > k_1/k \right\},$$

$$c^{-1} = \inf \left\{ \frac{\sigma}{\|\mathbf{x}\|}; \sigma, \mathbf{x}, f \text{ s.t. } \hat{p} < k_2/k \right\},$$

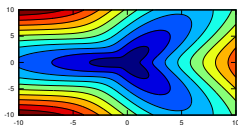
which provide assumptions 2 and 5 with  $b < c$

Assumptions 2 and 5 then imply  $b < b'$  and  $c' < c$

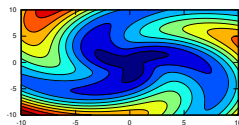
# Non quasi-convex functions



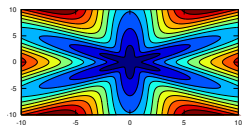
(a)



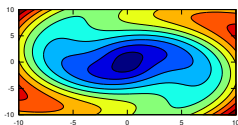
(b)



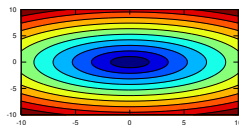
(c)



(d)



(e)



(f)